

Security Challenges in Virtualized Environments

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Confidence 2008,
Krakow, Poland, May 15th, 2008

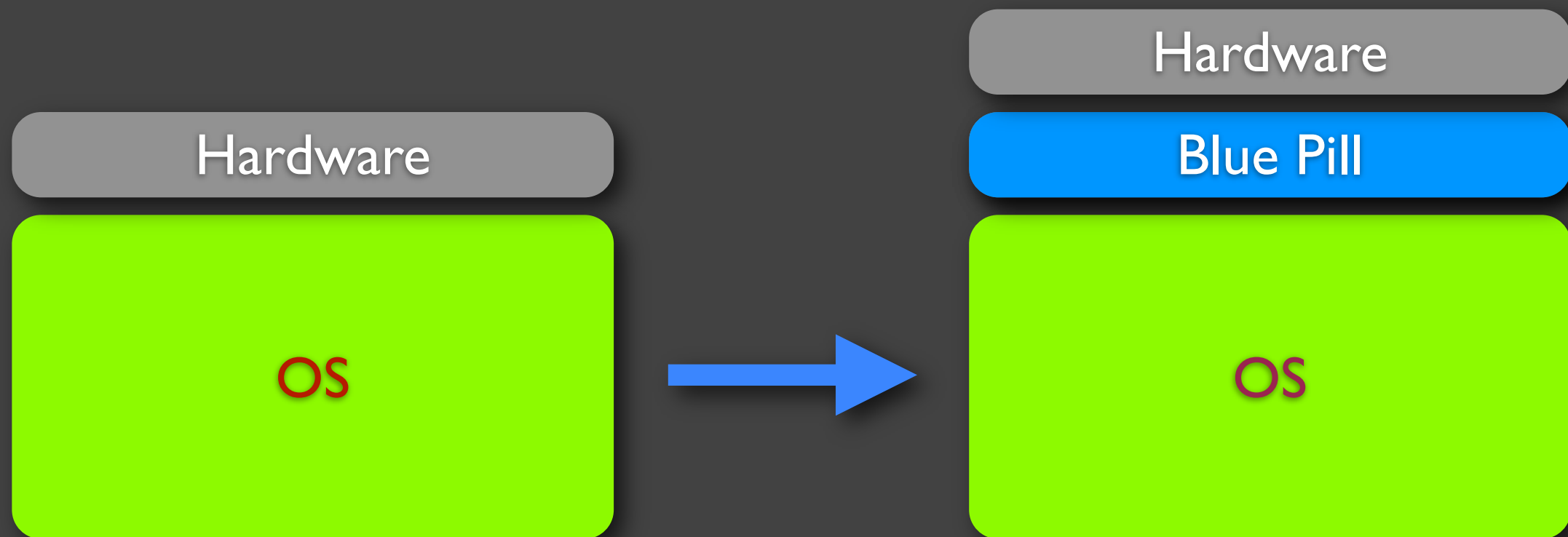
1 Virtualization-based **MALWARE**

2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



Virtualization-based MALWARE



AMD-V
Intel VT-x

Blue Pill Characteristics

NO HOOKS!



Cannot be detected using any integrity scanner

On the fly installation



No boot/BIOS/etc modifications necessary

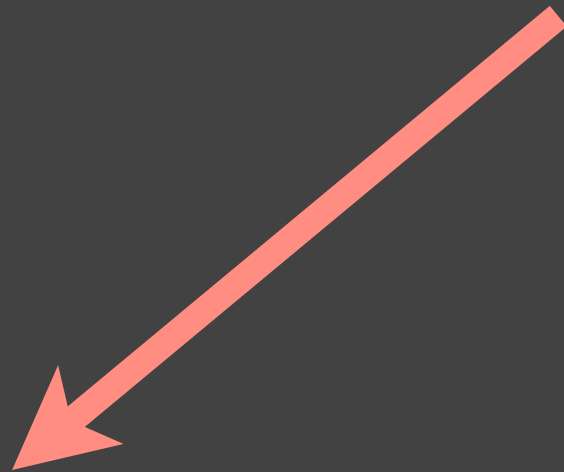
No I/O virtualization



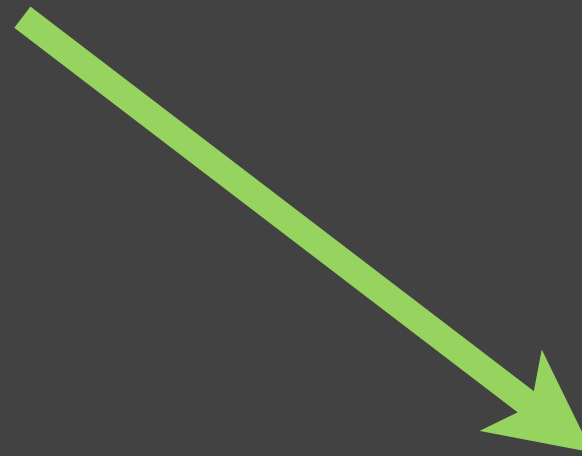
Negligible performance impact (your brand new 3D card will still work!)

Blue Pill detection

Blue Pill detection



Detecting a VMM



Detecting
virtualization based
malware

VMM detection

Direct timing analysis

Guest time virtualization

HPET timers

Blue Chicken

CPU specific behavior

TLB profiling

VMM detection?

- Everything is going to be virtualized!
- Thus the information that “there is a hypervisor in the system”...
- ...would be pretty much useless...

Detecting virtualized
malware?

No Hooks!

Search for code

Detect activity
(e.g. network packets)

Heuristics

By Pattern

- Stealth by Design concept
- Covert channels

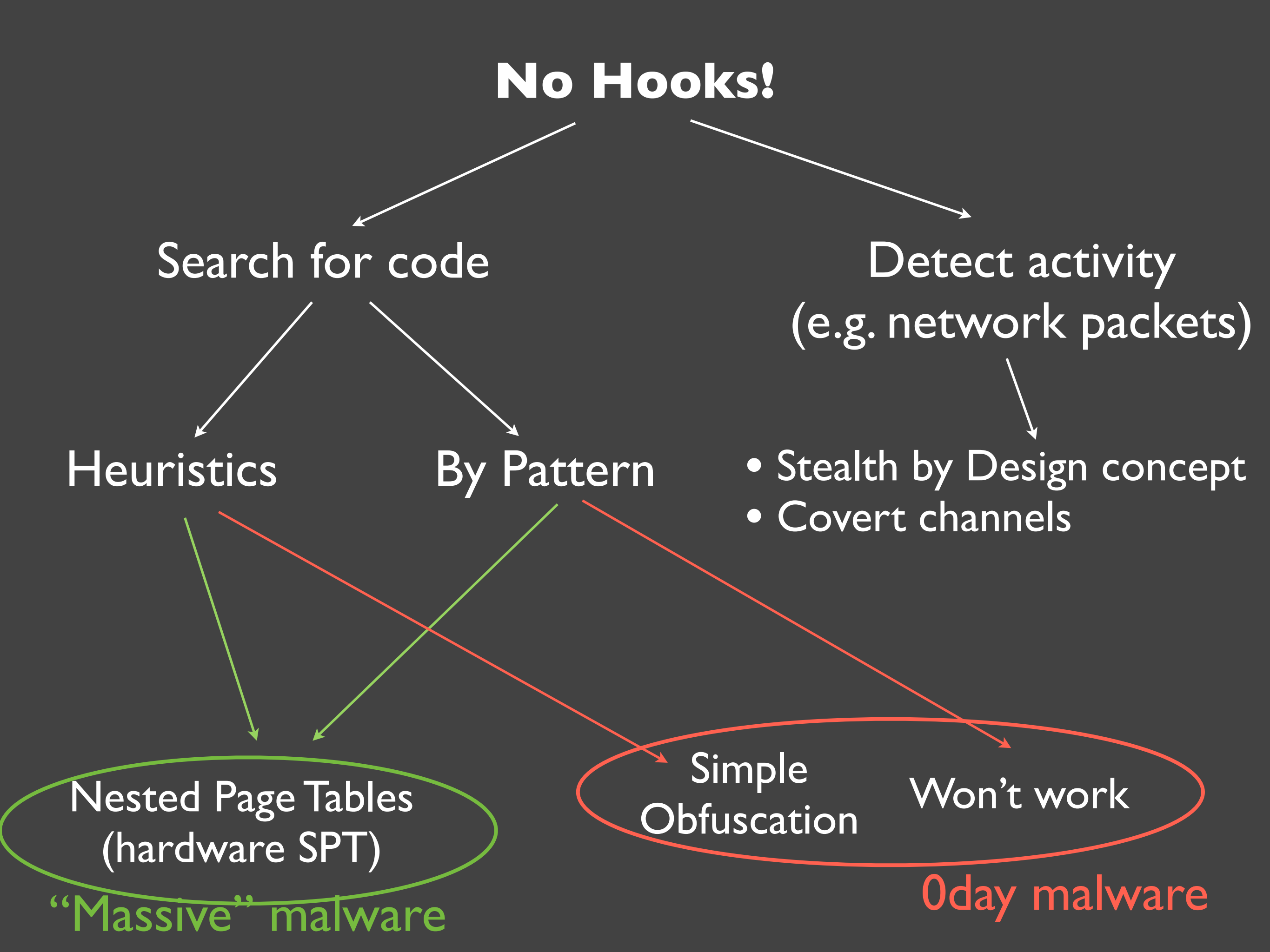
Nested Page Tables
(hardware SPT)

“Massive” malware

Simple
Obfuscation

Won't work

0day malware



The whole big deal about Blue Pill is:
NO HOOKS in the system!

Blue Pill prevention

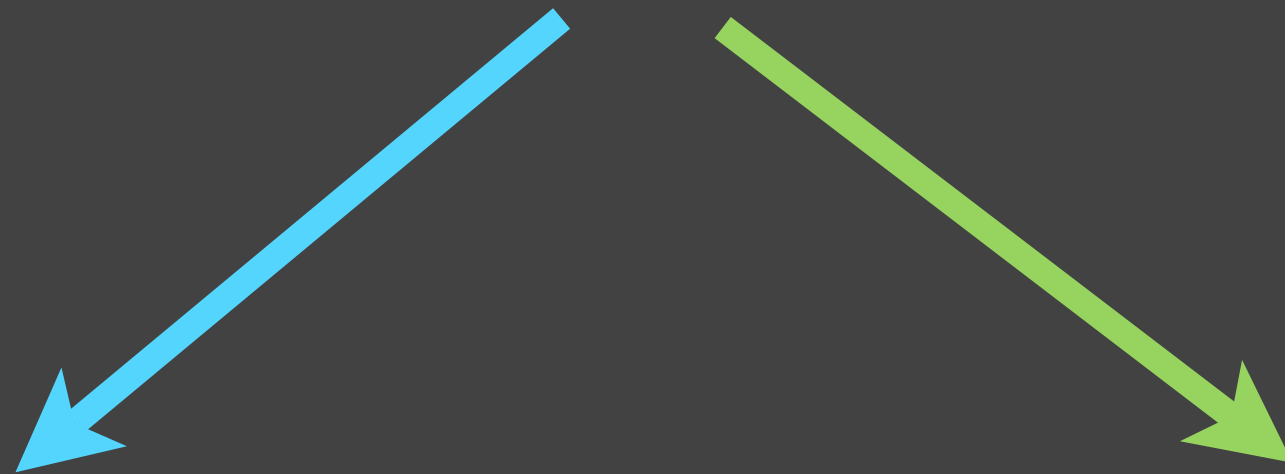


Disable virtualization?

How about also disabling your network card so you never got infected from the Internet?

Install a **trusted**
hypervisor first?

Installing trusted hypervisor



Static Root of Trust
Measurement

BIOS > MBR > VMM
e.g. MS Bitlocker

Dynamic Root of Trust
Measurement

SENTER (Intel TXT)
SKINIT (AMD SVM)

Trusted vs. Secure?

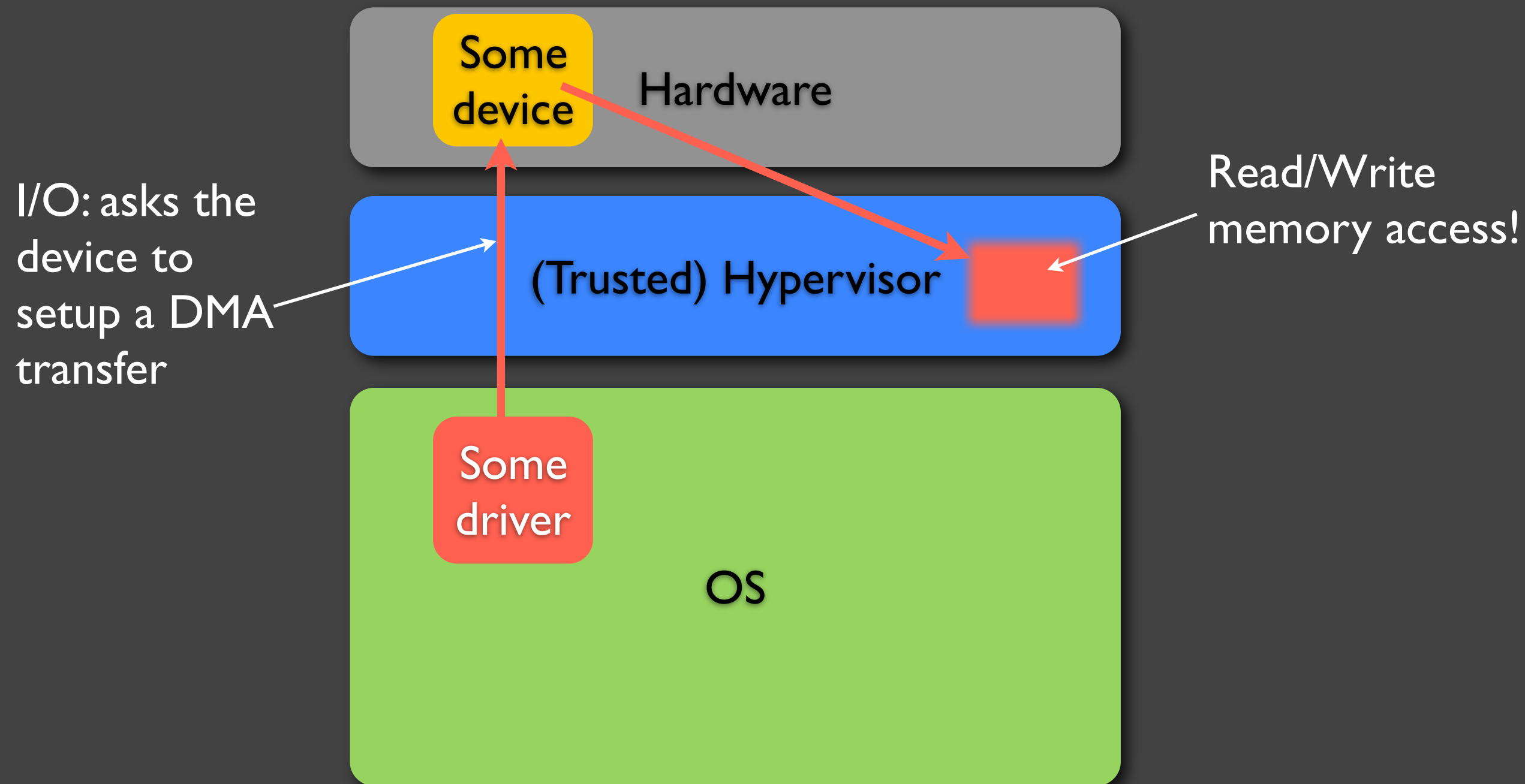
- SRTM and DRTM only assures that what we **load** is **trusted**...
- ...at the moment of **loading**!
- 3 sec later... it could be exploited and get compromised!

Trusted \neq Secure

*known,
not compromised*

no bugs

E.g. #1: The famous DMA problem



IOMMU

- Solution to the problem of “DMA attacks”
- Intel calls it: VT-d
- Not much PC hardware supports it yet
 - Expected to change soon
- No THIN HYPERVISORS without IOMMU!

Other problems with VMMs?
Stay tuned...

All in all:

it's not trivial to have a **trusted** & **secure** hypervisor..

... but this is the proper way to go!

Virtualization-based **MALWARE**

2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



Using Virtual Machines for ISOLATION

Originally **ISOLATION** was supposed to be provided by
Operating Systems...

- Separate processes/address spaces,
- User accounts & ACLs...

But in practice current OSes simply
fail at providing isolation!

Why OSes fail?

- Kernel bugs!
- Kernel bugs!!
- Kernel bugs!!!
- Bad design, e.g.:
 - XP and “all runs as admin” assumption
 - Vista’s UAC assumes admin rights should be granted to every installer program!

VMMs for the rescue!

trusted & secure hypervisor

Vista
(work projects)

Linux + Firefox
("random"
surfing)

Linux + Firefox
(online banking)

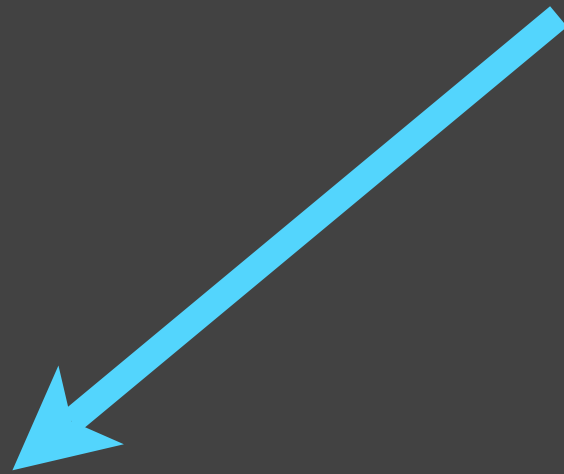
MacOSX
("home", e.g.
pics, music, etc)

Challenges

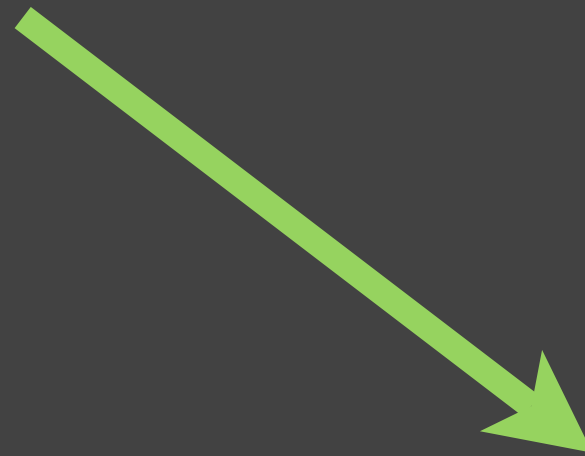
- Performance
- Why is VMM/hypervisor going to be more secure than OS's kernel?

VMM bugs?

VMM Bugs



Bugs in hypervisors



Bugs in additional
infrastructure

E.g. #1: CVE-2007-4496

- VMWare ESX 3.0.1
 - <http://www.vmware.com/support/vi3/doc/esx-8258730-patch.html>
- Found by Rafal Wojtczuk (McAfee)
- September 2007
- Guest OS can cause memory corruption on the host and *potentially* allow for arbitrary code execution on the host

E.g. #2: CVE-2007-0948

- Microsoft Virtual Server 2005 R2
 - <http://www.microsoft.com/technet/security/bulletin/ms07-049.msp>
- Found by Rafal Wojtczuk (McAfee)
- August 2007
- Heap-based buffer overflow allows guest OS to execute arbitrary code on the host OS

E.g. #3: CVE-2007-4993

- Xen 3.0.3
 - http://bugzilla.xensource.com/bugzilla/show_bug.cgi?id=1068
- Found by Joris van Rantwijk
- September 2007
- By crafting a grub.conf file, the root user in a guest domain can trigger execution of arbitrary Python code in domain 0.

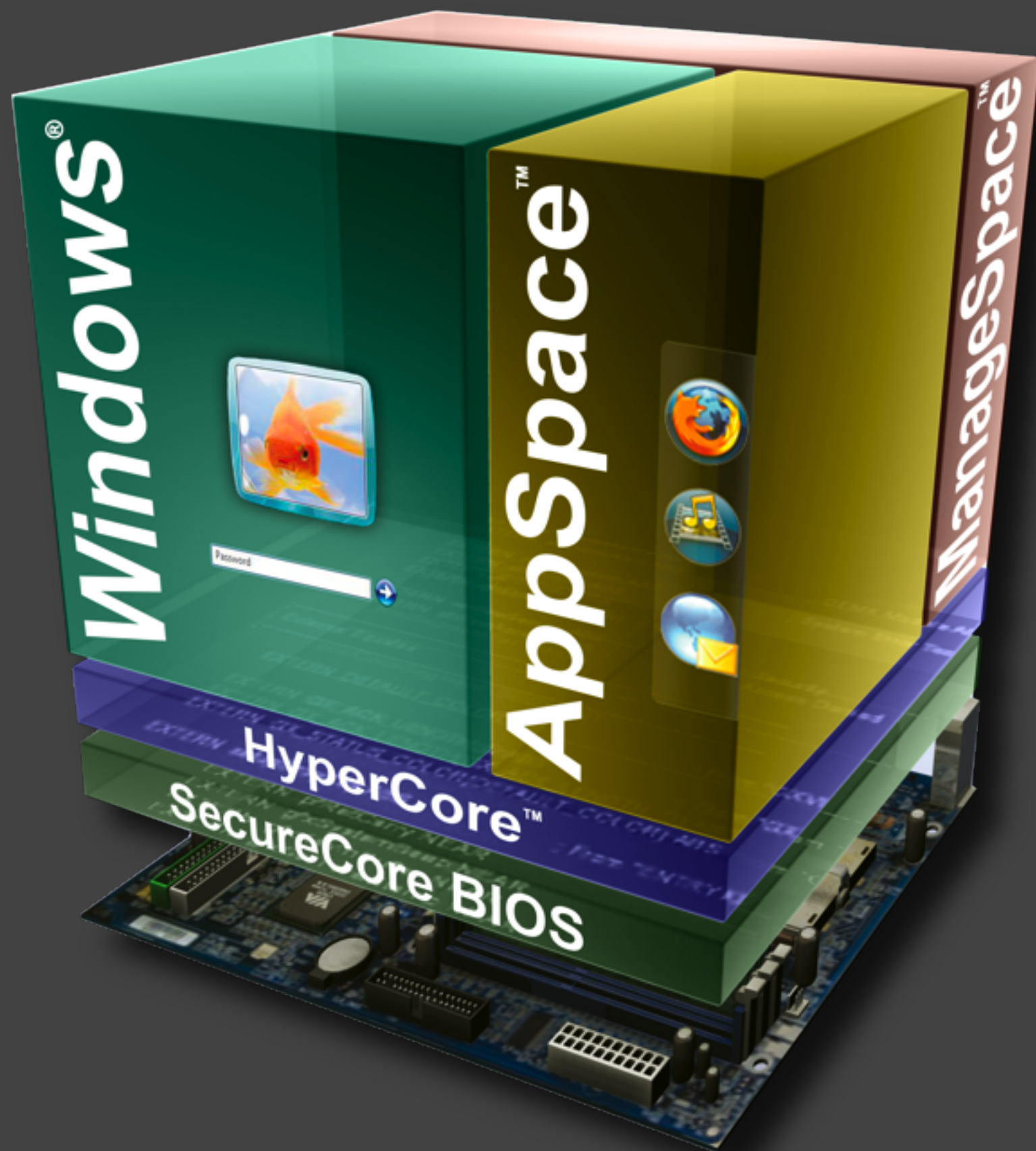
E.g. #4: Various Bugs

- Paper by Tavis Ormandy (Google)
 - <http://taviso.decsystem.org/virtsec.pdf>
- April 2007
- Disclosed bugs in VMWare, XEN, Bochs, Virtual PC, Prallels
- A simple fuzzers for:
 - Instruction parsing by VMMs
 - I/O device emulation by VMMs

As you see, current VMMs are far from being flawless...

To make VMMs more secure we need to keep them
ultra-thin and small!

Phoenix HyperSpace



Search

02:39 pm

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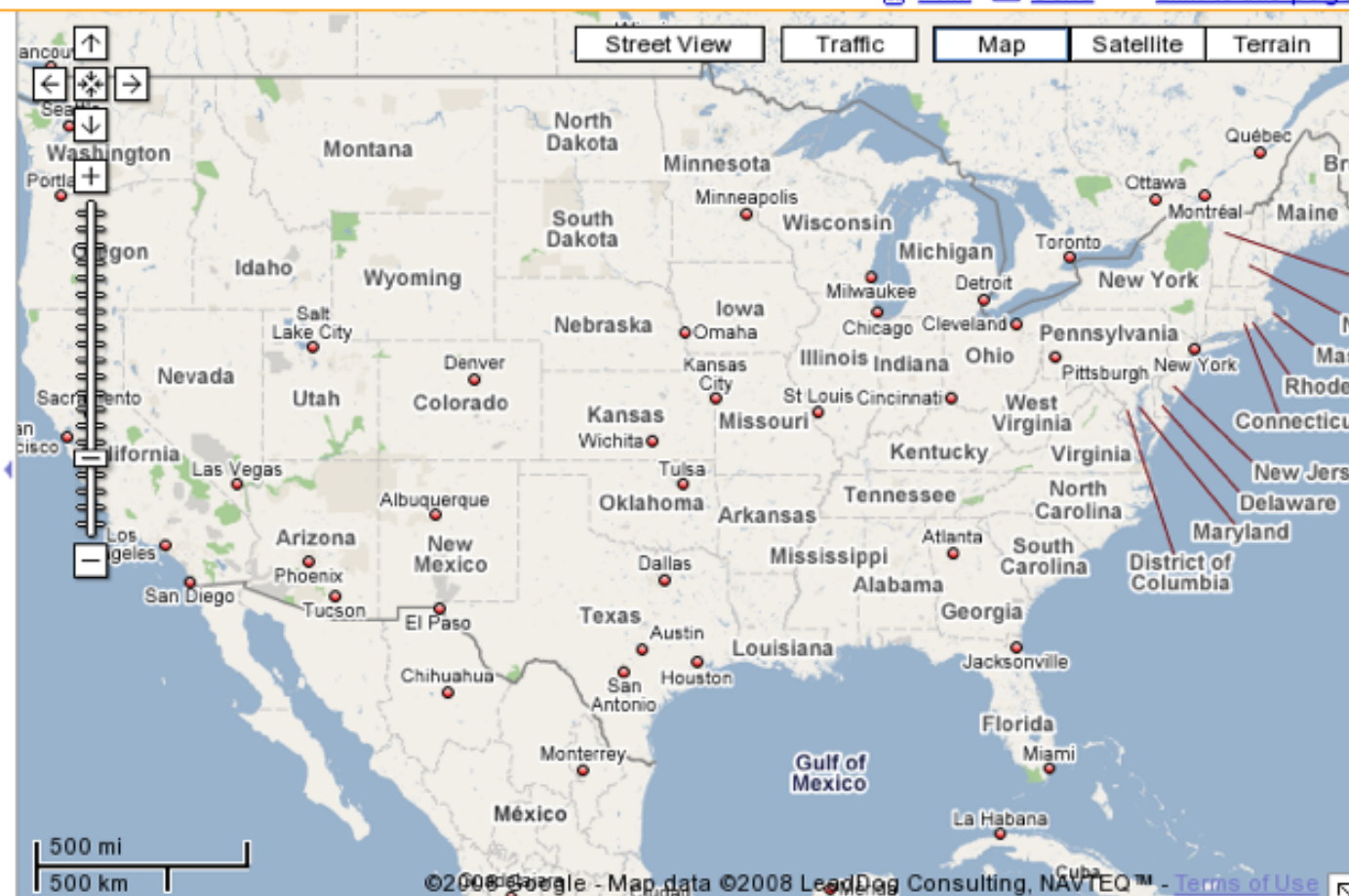
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Done

Two Kennedys endor...

MONDAY, Jan 20

5 U.S. troops killed...

MONDAY, Jan 20

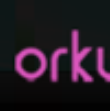
Burglar steals bishop...

MONDAY, Jan 20

Militants release Pak...

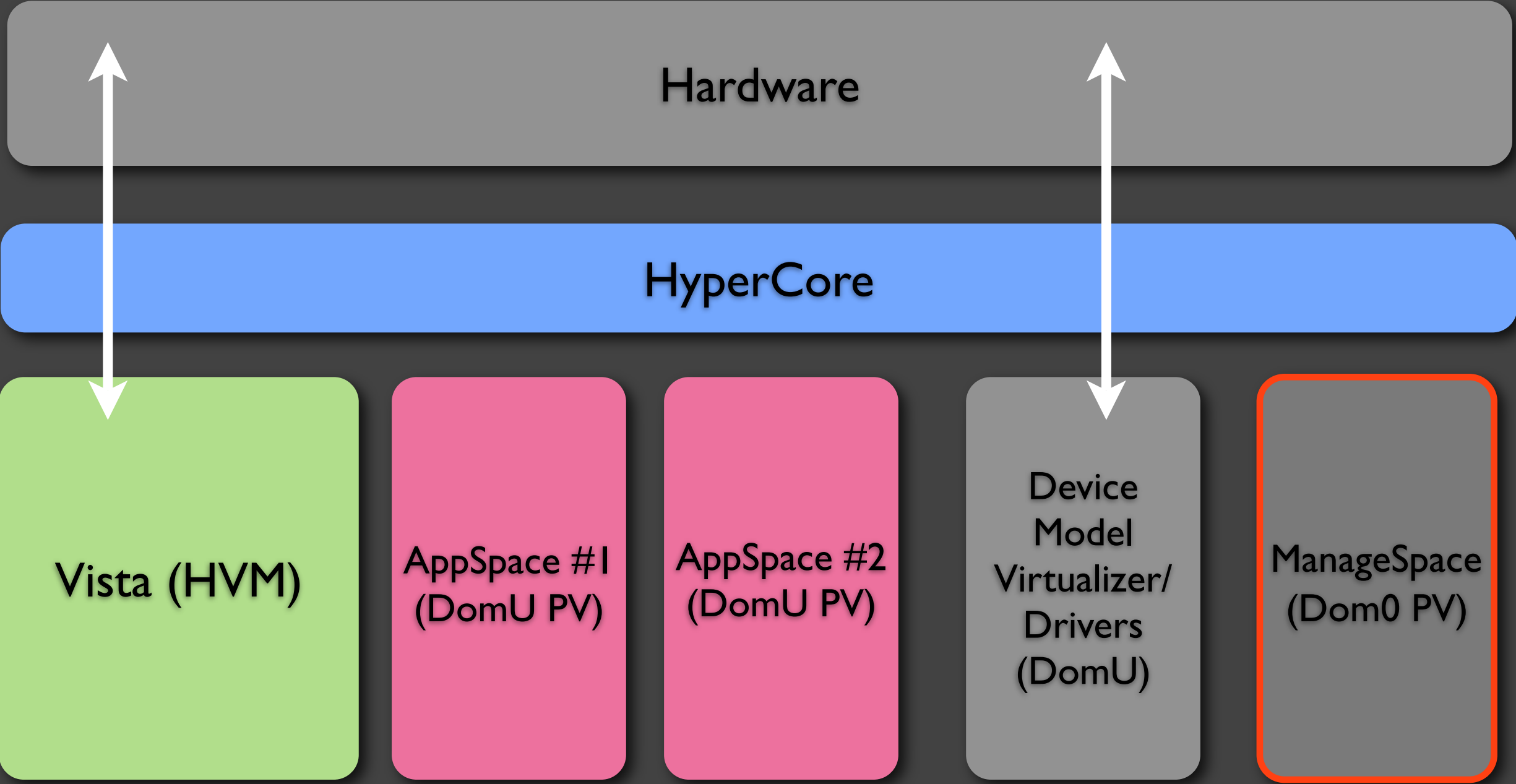
MONDAY, Jan 20

9-12



2:39 PM

HyperCore:
the type I hypervisor used for HyperSpace



The HyperCore

- Targets desktop/laptop systems
- Guest OS execute at near-native performance (including fancy graphics)
- Support for full ACPI (Power Management)
- Integrity: loaded via SecureCore BIOS (Static Root of Trust Measurement)
- Very thin - easy to audit!

Speeding things up

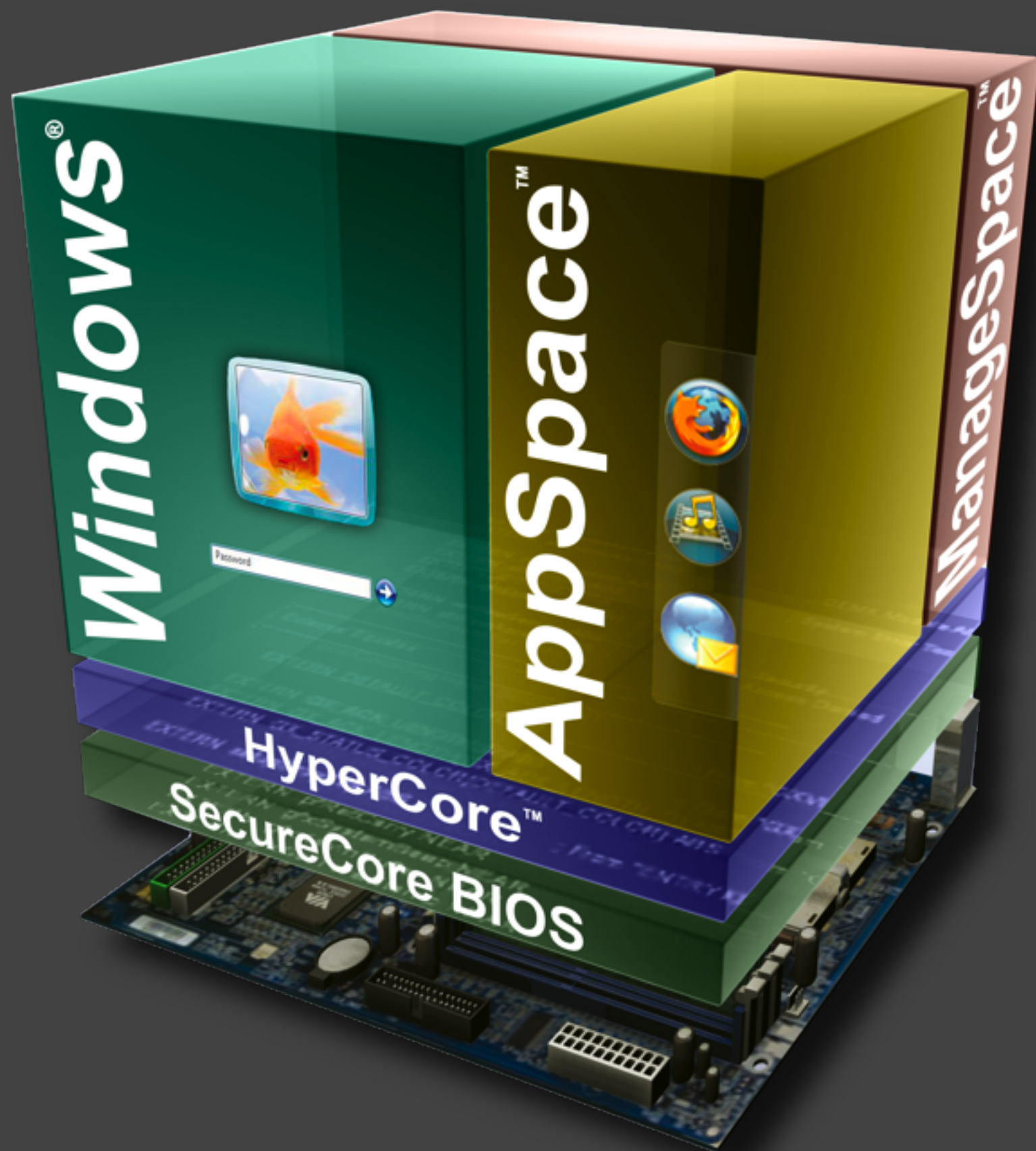
- Pass through for most devices
- SPT: I-I mapping for most pages for the Primary OS

Power Management

- ACPI tables exposed to the Primary OS, so that the overall power performance is optimized
- Efficient intercepts for power management control

Integrity

- Static RTM via Phoenix's SecureCore BIOS
- Dynamic RTM via Intel's TXT/AMD's SKINIT
- SMM-based watchdog for HyperCore code



1 Virtualization-based **MALWARE**

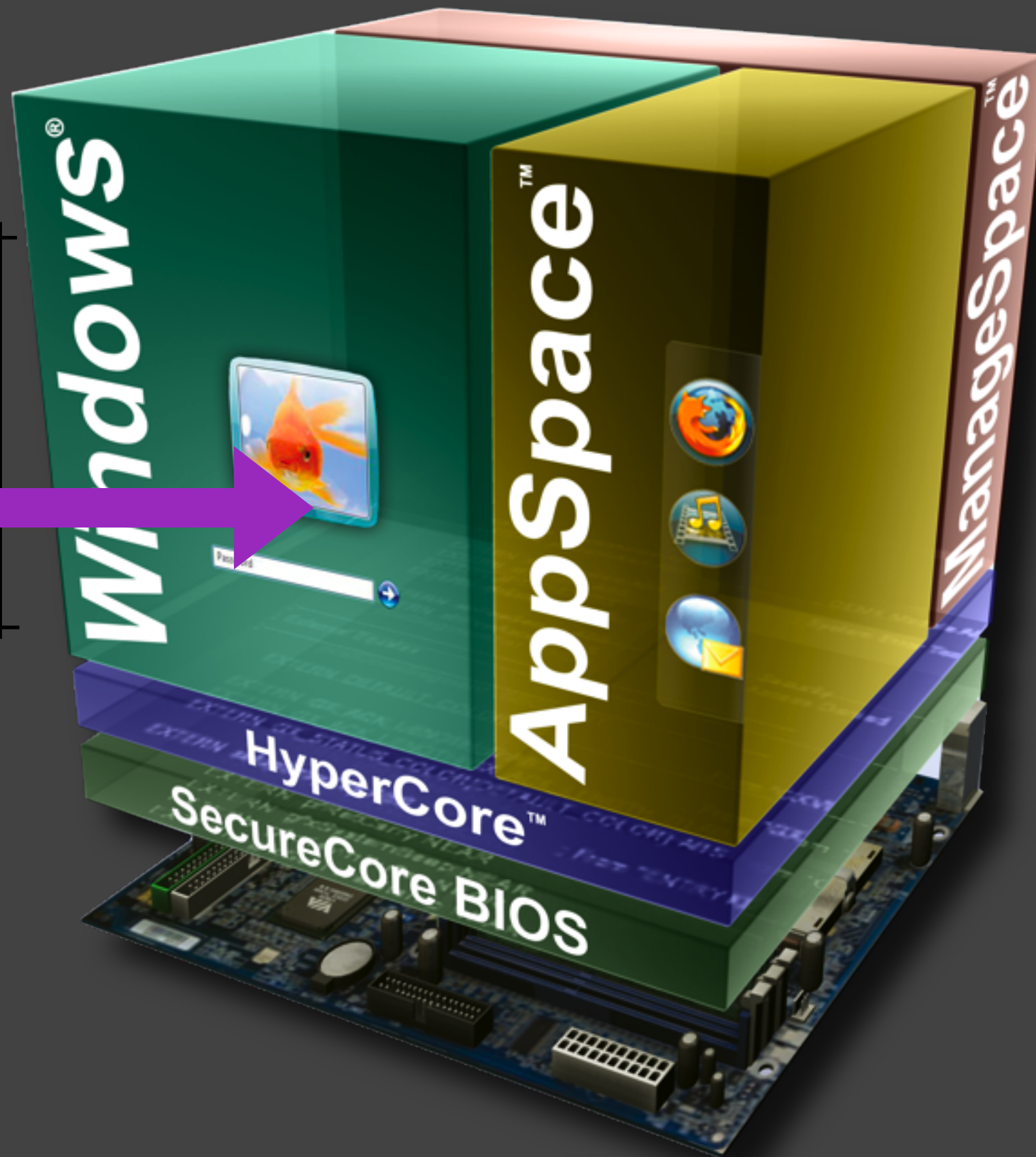
2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



NESTED virtualization

What if a
user wants
to run e.g.
Virtual PC
here?



Hypervisor (Primary)

VM₁

VM₂ (Nested
Hypervisor)

VM₃

VM₄

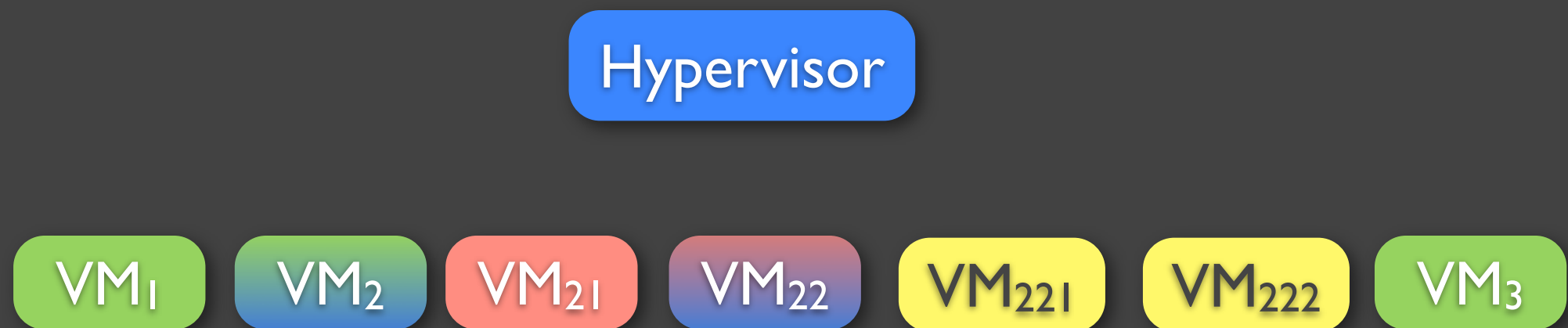
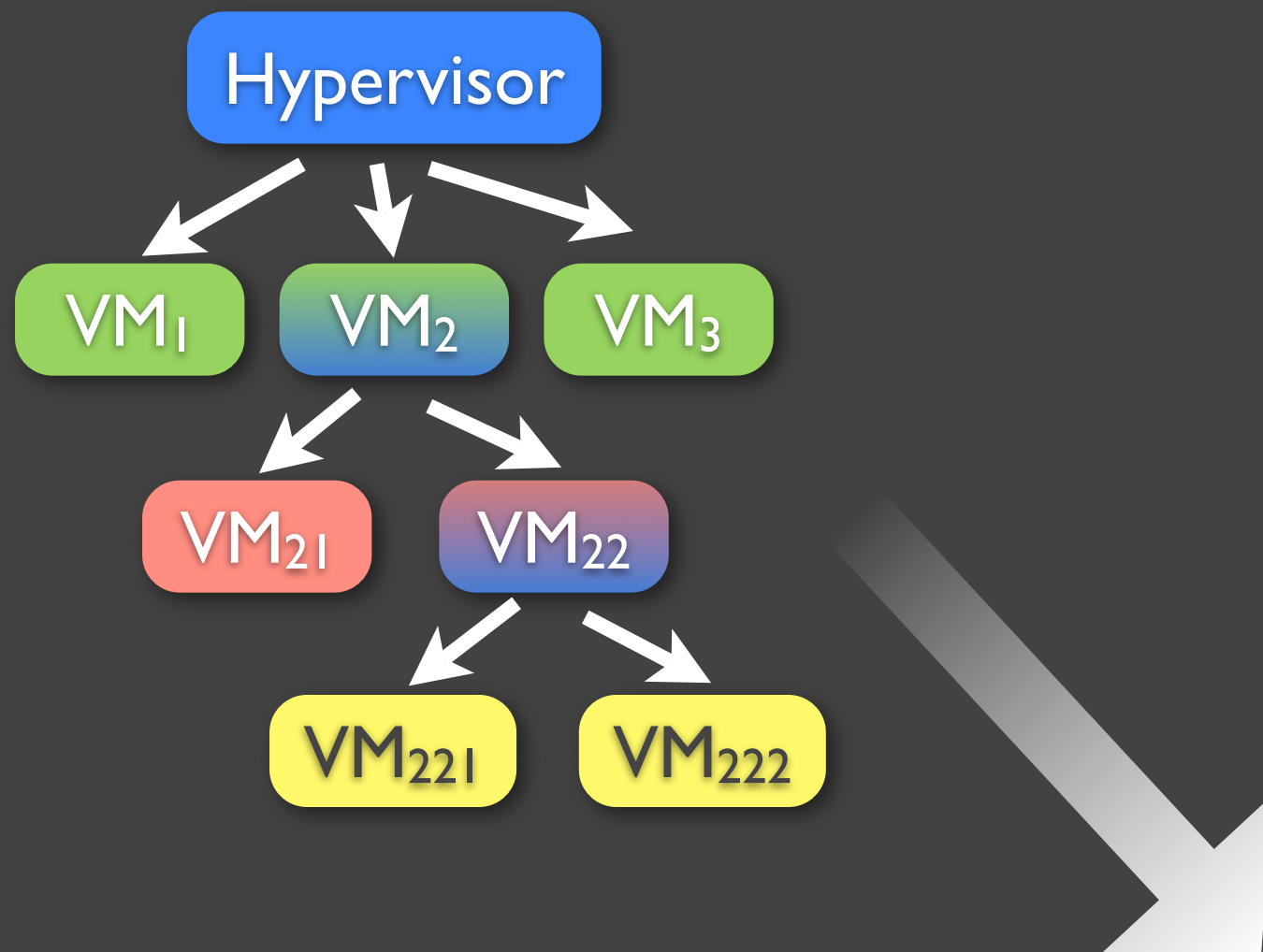
VM₂₁

VM₂₂

VM₂₂₁

VM₂₂₂

Idea of how to handle this situation...



Now, lets look at the actual details :)

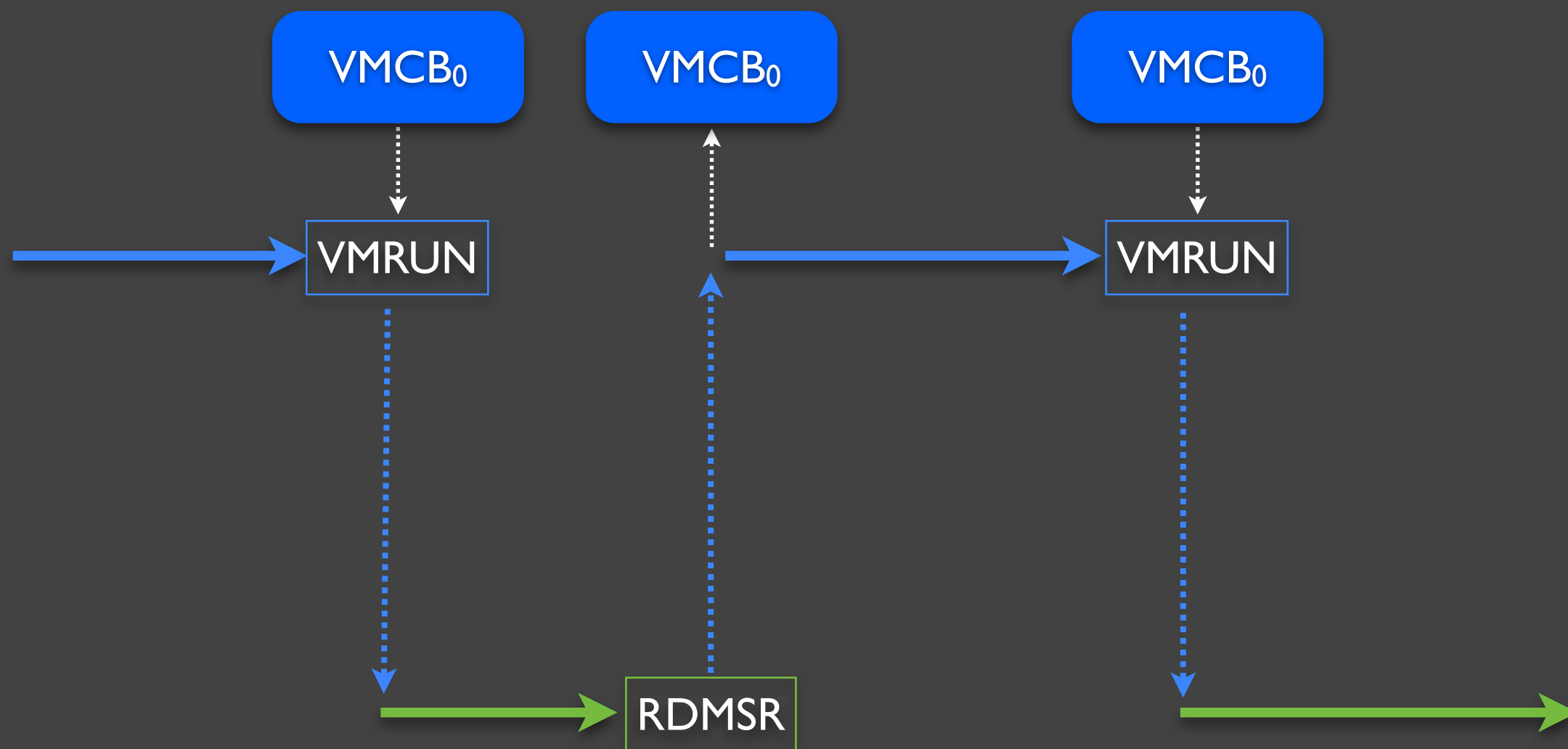
Let's start with AMD-V...

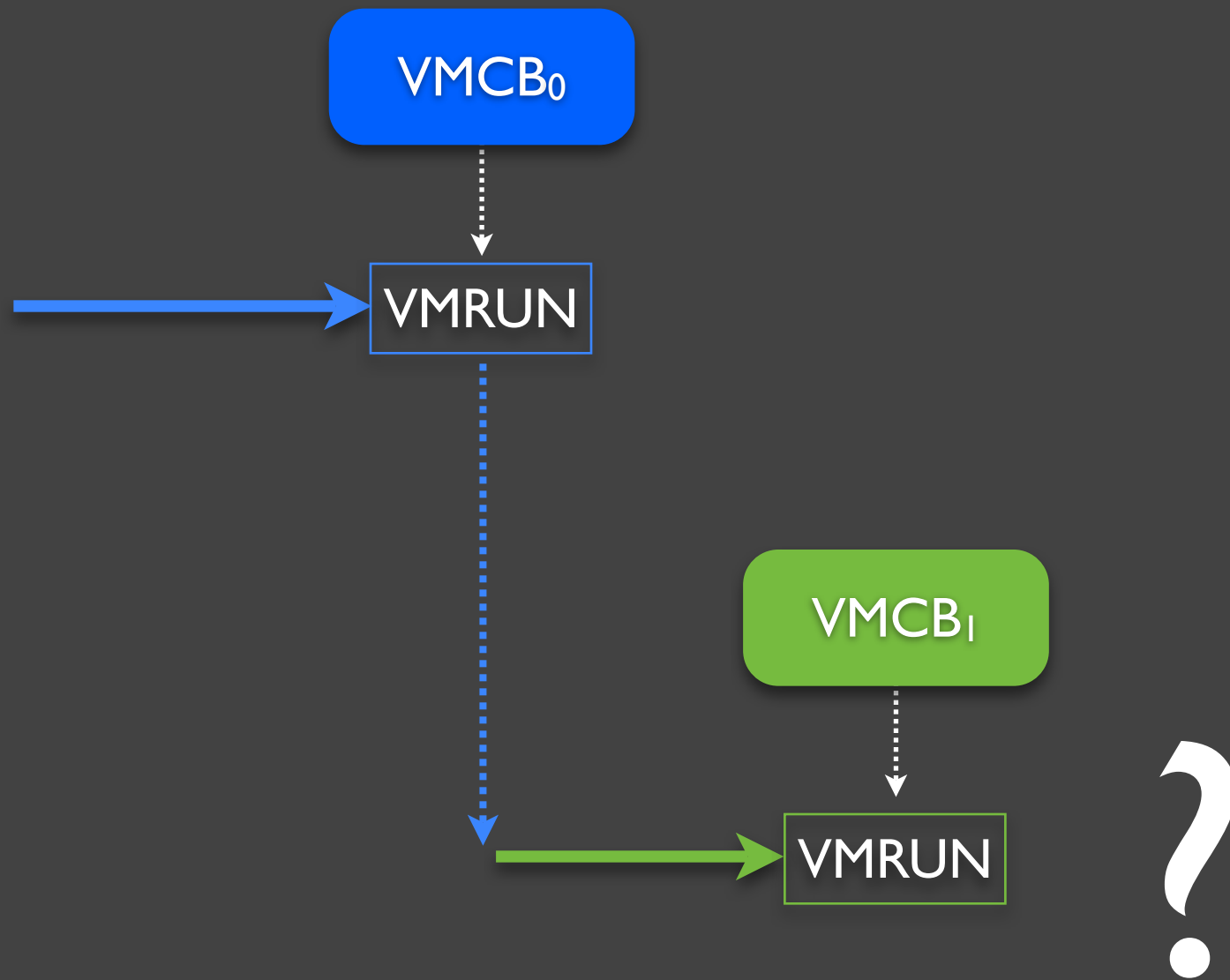
WARNING!!!

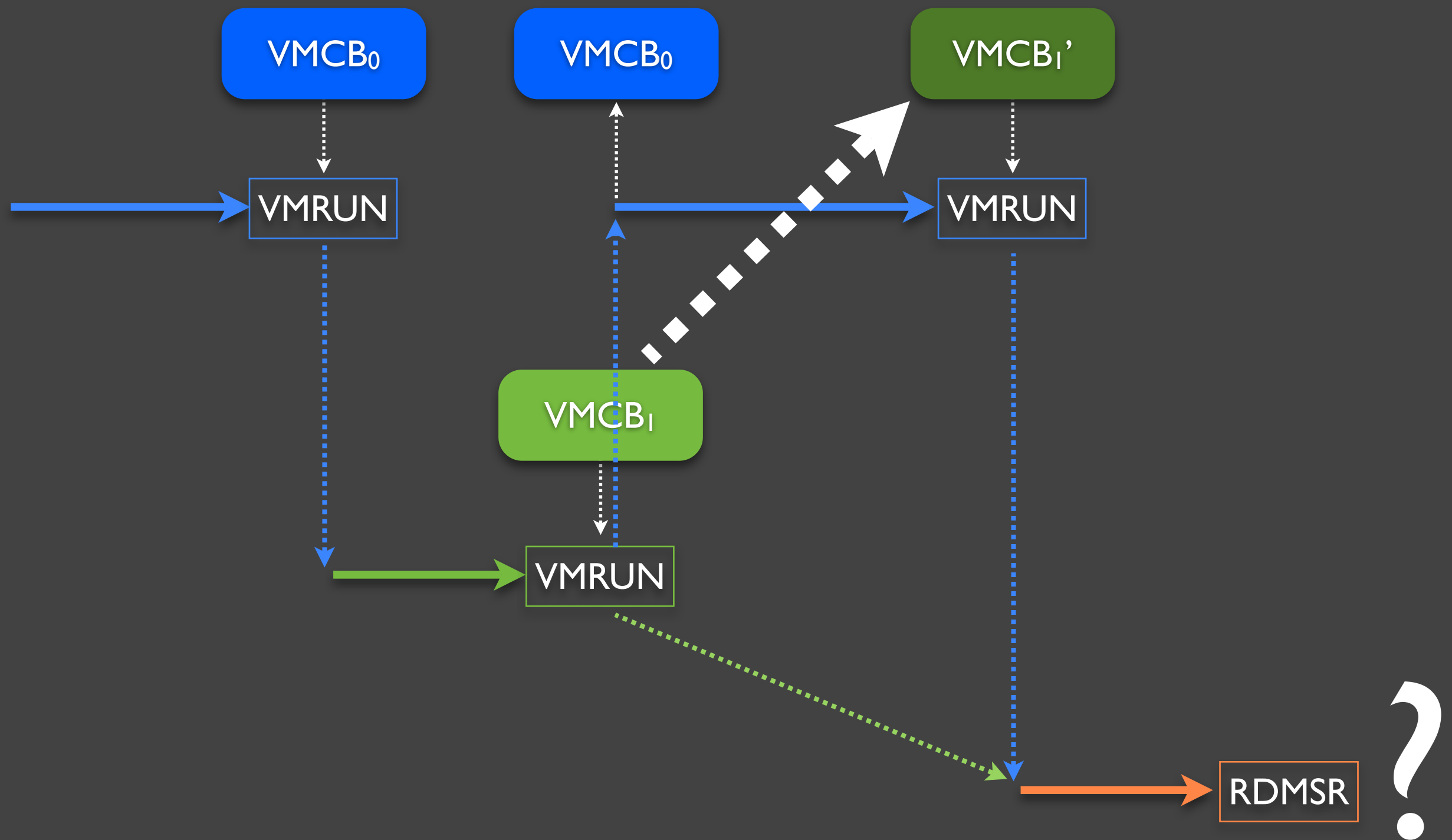
deep technical content follows

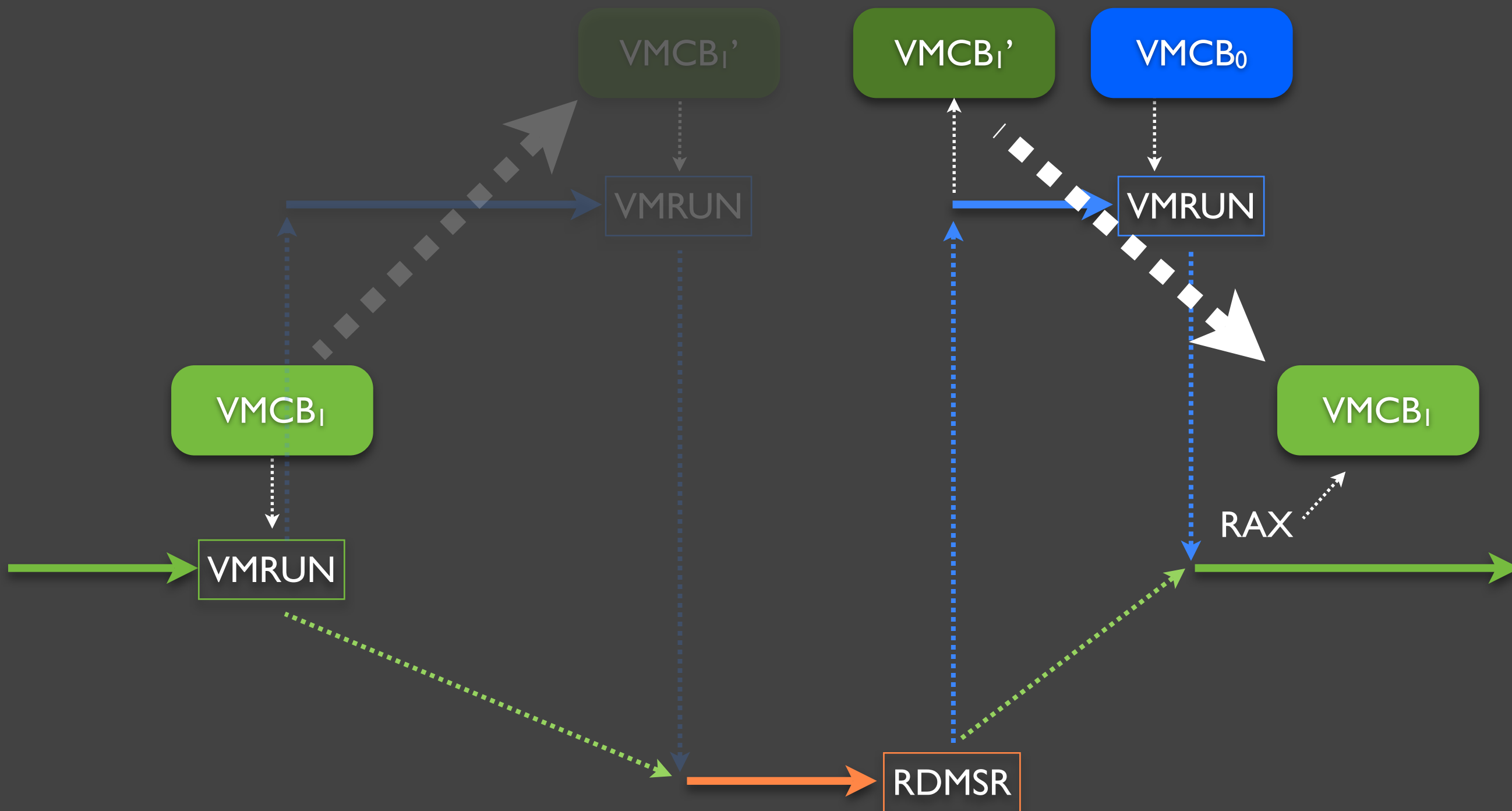
(next 14 slides)

don't worry if you get lost!

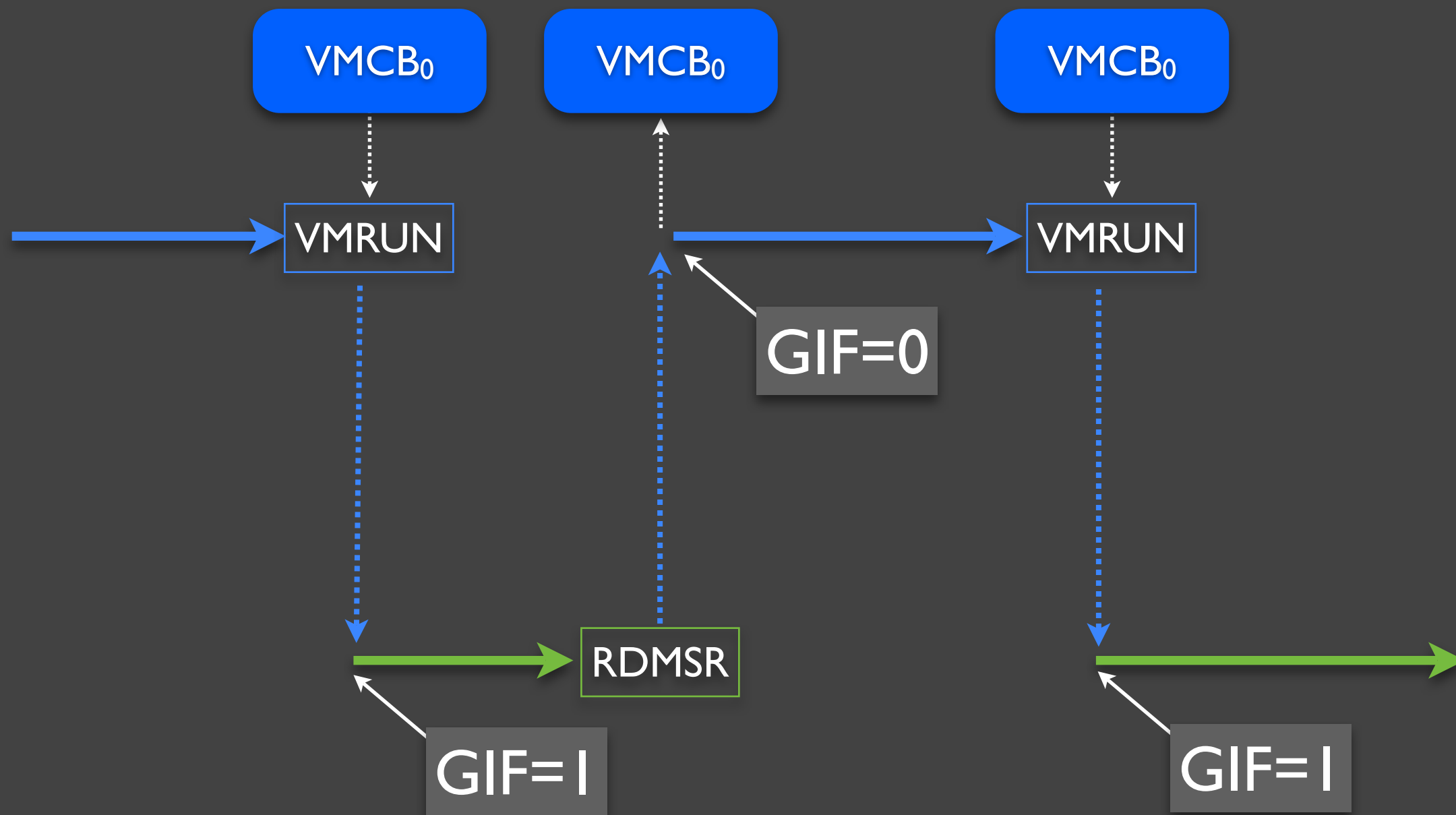


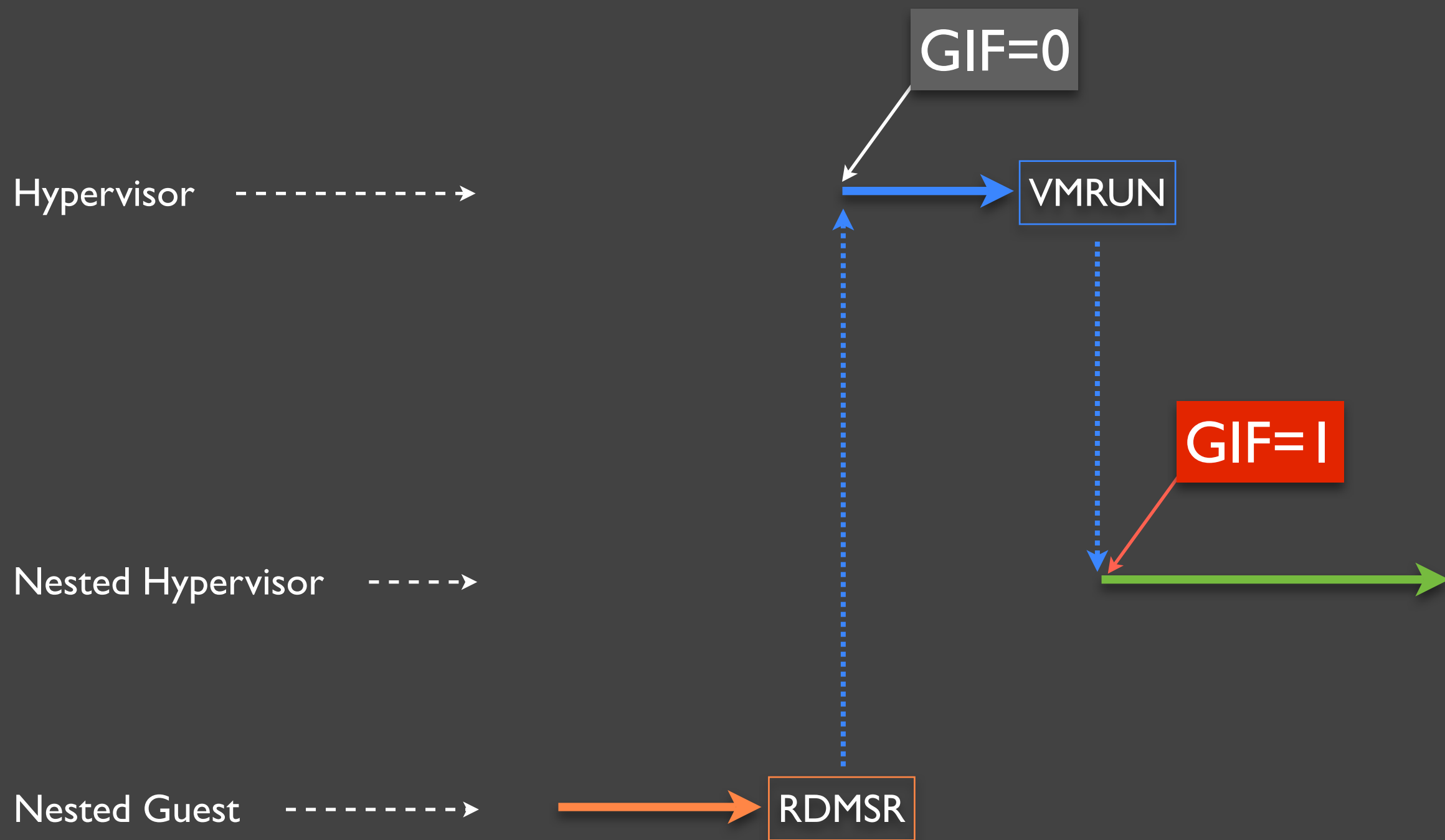






Looks convincing but we also need to take care about
some technical details, that are not trivial...





- Hypervisors expect to have GIF=1 when VMEXIT occurs...
- They might not be prepared to handle interrupts just after VMEXIT from guests!
- ... but when we resume the nested hypervisor CPU sets GIF=1, because we do this via VMRUN, not VMEXIT...

Getting around the “GIF Problem”

- We need to “emulate” that GIF is 0 for the nested hypervisor
- We stop this emulation when:
 - The nested hypervisor executes STGI
 - The nested hypervisor executes VMRUN
- How do we emulate it?

GIF0 emulation

- $\text{VMCB}_i.\text{V_INTR_MASKING} = 1$
- Host's $\text{RFLAGS.IF} = 0$
- Intercept NMI, SMI, INIT, #DB and held (i.e. record and reinject) or discard until we stop the emulation

Additional details

- Need to also intercept VMLOAD/VMSAVE
- Need to virtualize VM_HSAVE_PA
- ASID conflicts

Hypervisor:ASID = 0

Conflicting ASIDs!

Nested Hypervisor:ASID = 1
(but thinks that has ASID = 0)

Nested Guest:ASID = 1
(assigned by the nested hypervisor)

But we can always reassign the ASID in the VMCB “prim”
that we use to run the nested guest.

Performance Impact

- One additional #VMEXIT on every #VMEXIT that would occur in a non-nested scenario
- One additional #VMEXIT when the nested hypervisor executes: STGI, CLGI, VMLOAD, VMSAVE
- Lots of space for optimization though

Lost already? ;)

Don't worry! The main message is...

This can be done!
&
It works!

```
Administrator: Command Prompt
C:\tmp\nbp-0.30>\tools\w2k_load.exe bin\amd64\newbp.sys

// w2k_load.exe
// SBS Windows 2000 Driver Loader V1.00
// 08-27-2000 Sven B. Schreiber
// sbs@orgon.com

Loading "bin\amd64\newbp.sys" ... OK




C:\tmp\nbp-0.30>\tmp\bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp\nbp-0.30>"\Program Files (x86)\Microsoft Virtual PC\Virtual PC.exe"

C:\tmp\nbp-0.30>\tmp\bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp\nbp-0.30>
```

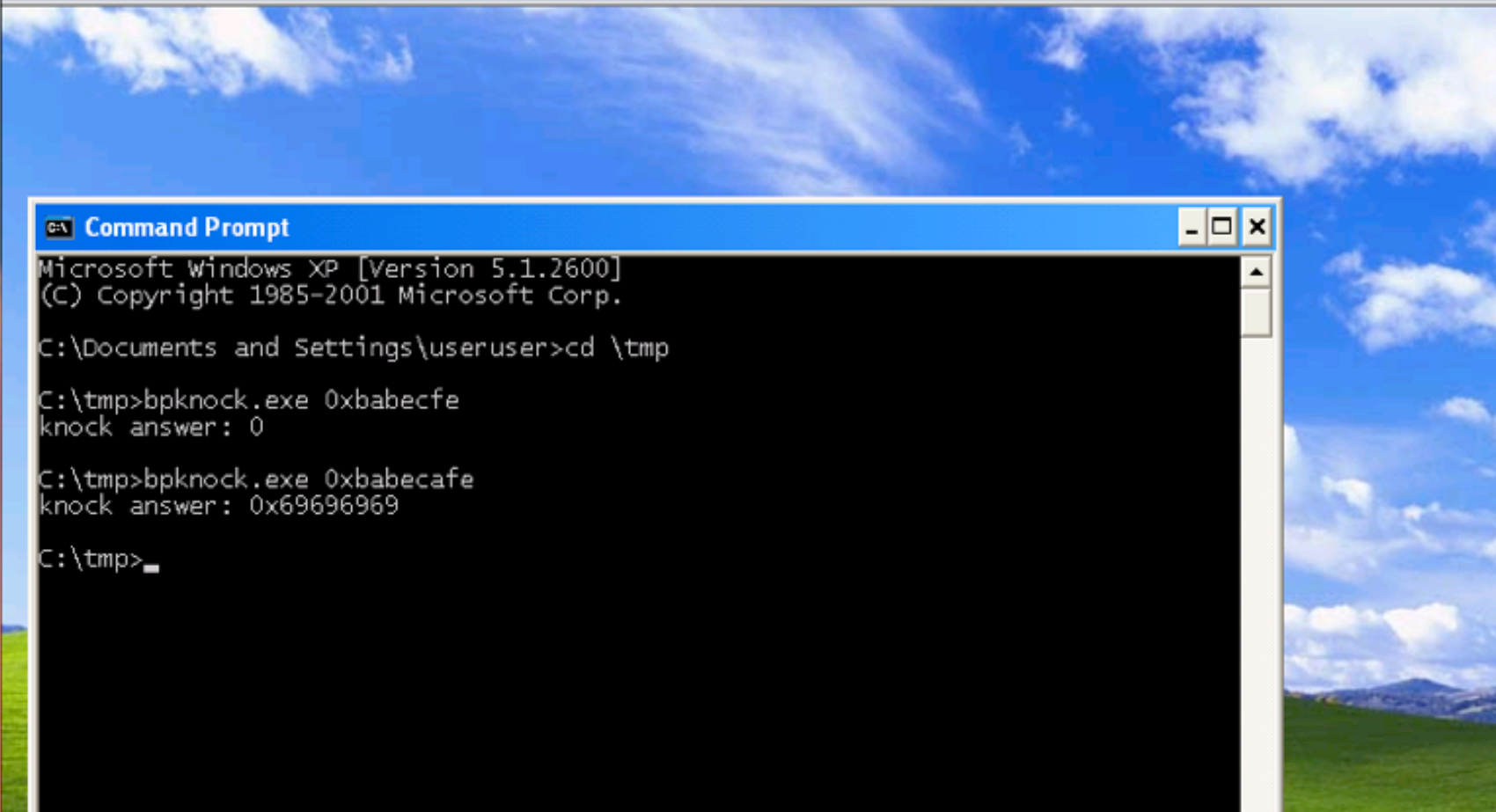
Virtual PC Console

File	Action	Help
	Linux install	Not running
	Vista	Not running
	Windows XP	Running

New...
Settings
Remove
Close...

Windows XP - Microsoft Virtual PC 2007

Action Edit CD Floppy Help



Command Prompt

```
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\useruser>cd \tmp

C:\tmp>bpknock.exe 0xbabecfe
knock answer: 0

C:\tmp>bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp>
```



<http://bluepillproject.org>

Intel VT-x

Nested virtualization on VT-x

- No GLF bit - no need to emulate “GLF0” for the nested hypervisor :)
- No Tagged TLB - No ASID conflicts :)
- However:
 - VMX instructions can take memory operands - need to use complex operand parser
 - No tagged TLB - potentially bigger performance impact

Nested VT-x: Status

- We have that working!
- The VT-x nesting code cannot be published though :(

Who else does Nested (hardware-based) Virtualization?

IBM z/VM hypervisor on IBM System z™ mainframe

“Running z/VM in a virtual machine (that is, z/VM as a guest of z/VM, also known as “second-level” z/VM) is functionally supported but is intended only for testing purposes for the second-level z/VM system and its guests (called “third-level” guests).”

-- <http://www.vm.ibm.com/pubs/hcsf8b22.pdf>

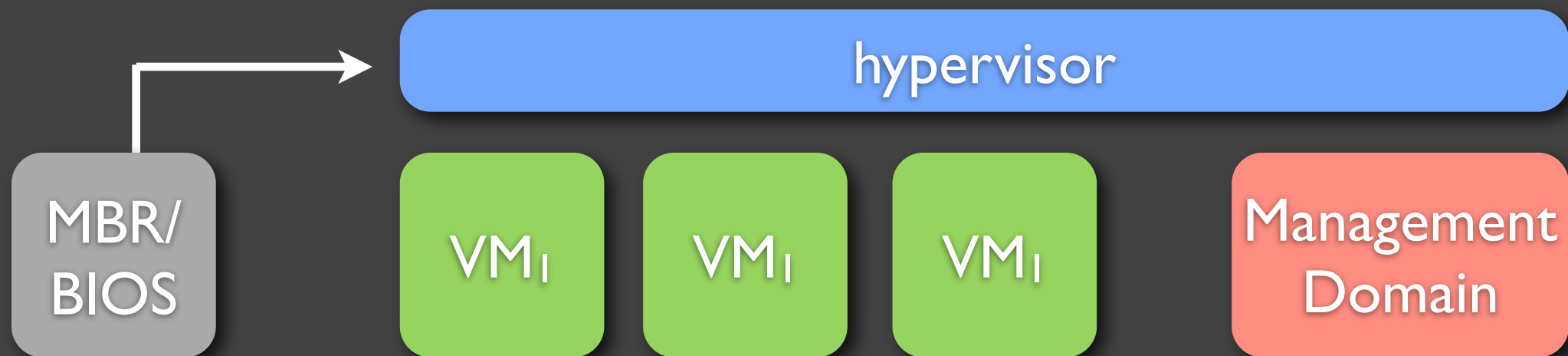


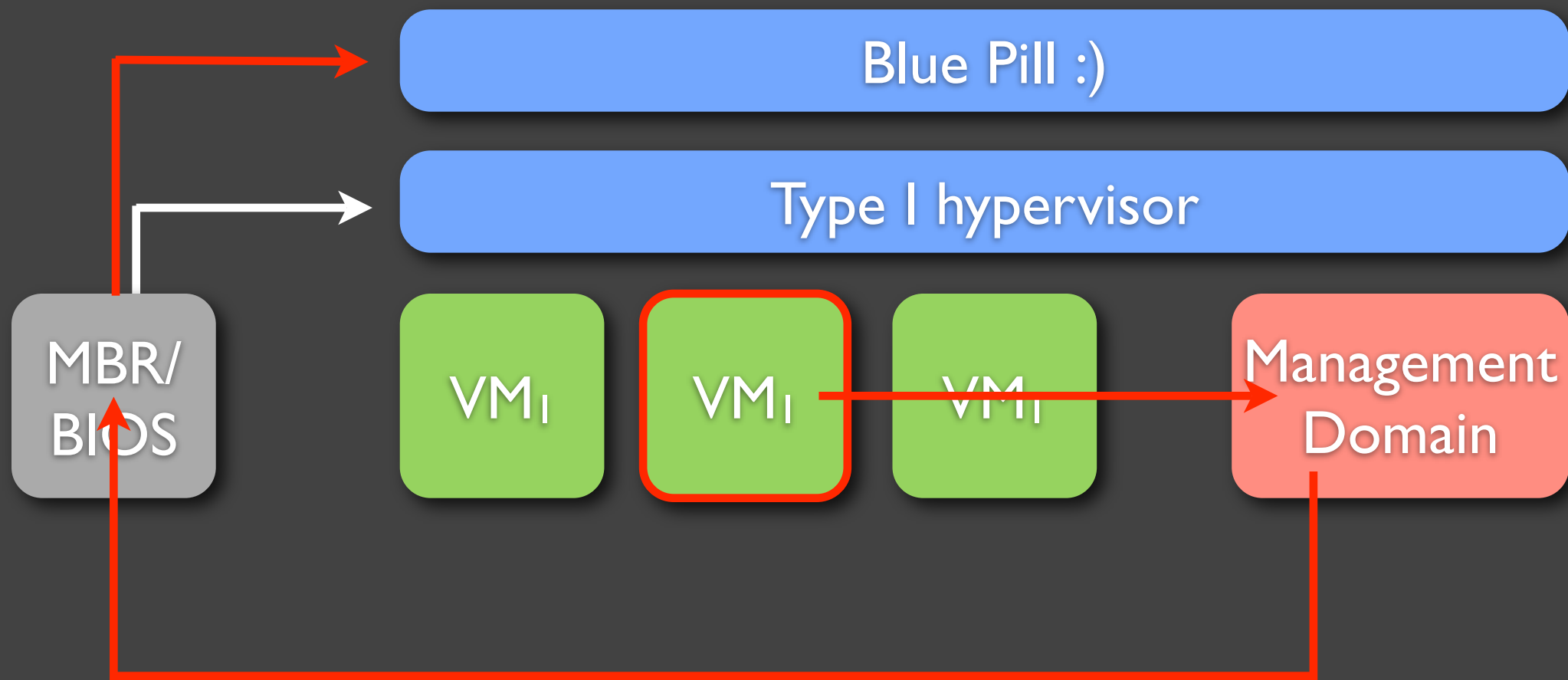
IBM System z10, source: ibm.com

Confusion

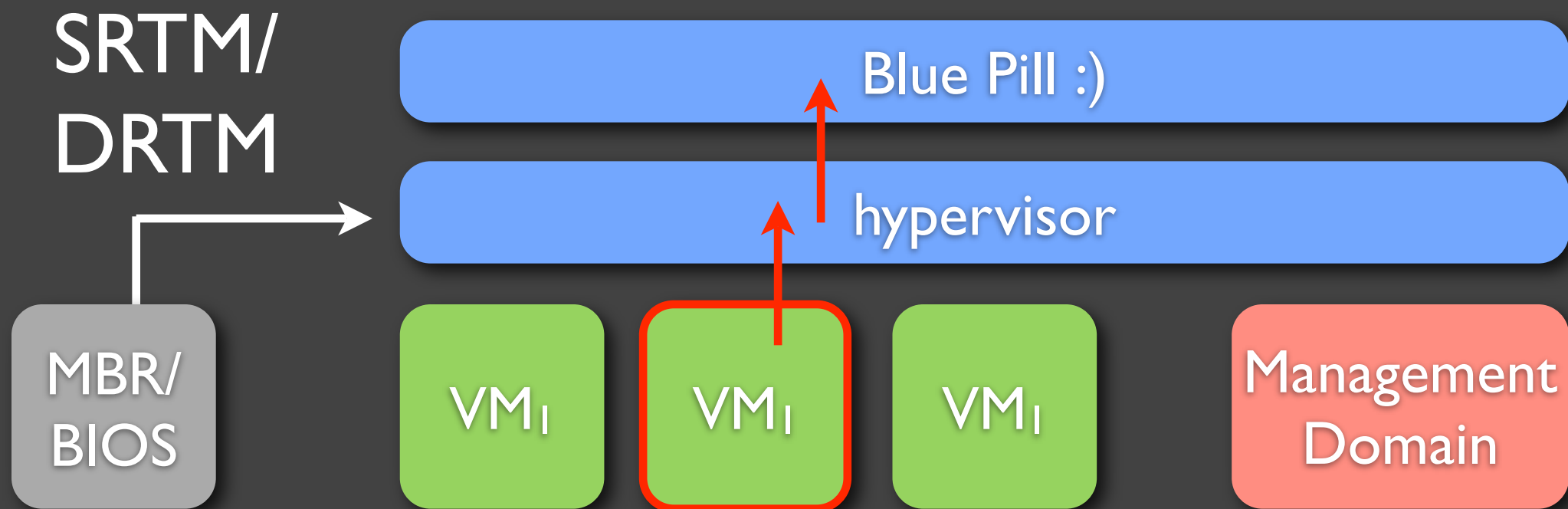
- AMD Nested Page Tables != Nested Virtualization!
- NPT is a hardware alternative to Shadow Page Tables (a good thing, BTW)
- NPT is also called: Rapid Virtualization Indexing

Nested Virtualization: Security Implications





Solution: ensure hypervisor integrity via SRTM or DRTM



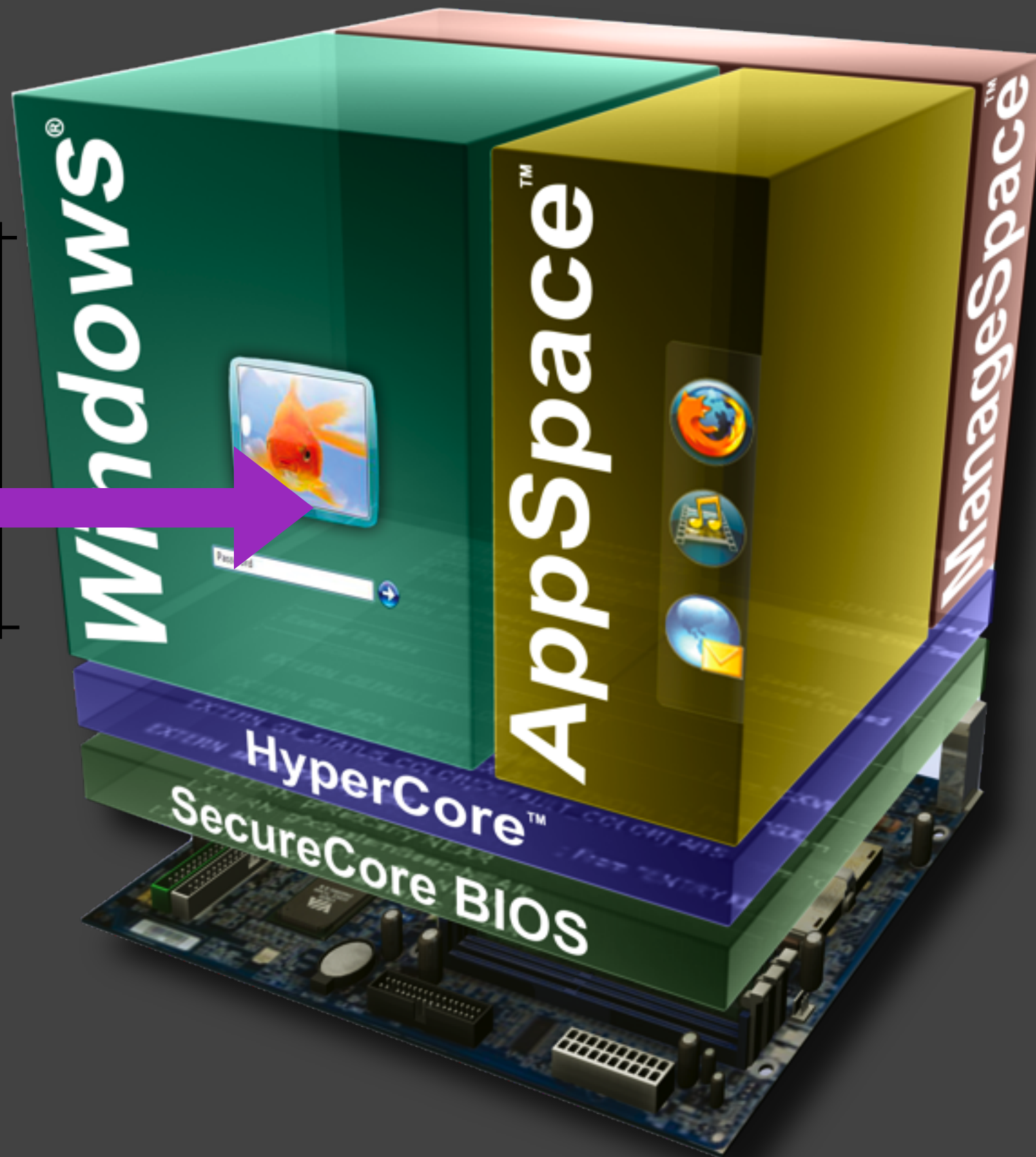
SRTM/DRTM do not protect the **already loaded hypervisor**, from being exploited **if it is buggy!**

Keep hypervisors very slim!
Do not put drivers there!

Nested Virtualization: Useful Applications

Remember Phoenix's HyperCore?

What if a
user wants
to run e.g.
Virtual PC
here?



Phoenix Technologies has supported the research on
nested hypervisors since Fall 2007

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Summary

- Virtualization technology could be used to improve security on desktop systems
- However there are non-trivial challenges in making this all working well...
- Virtualization is cool ;)

New stuff coming soon...

Summer 2008 in Las Vegas...

Invisible Things Lab
<http://invisiblethingslab.com>

Thank you!